## Galois connections for phylogenetic networks and their polytopes.

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## Trees



Episodic radiations in the fly tree of life, [Wiegmann et.al. PNAS 2011]

## Splits



1) Phylogenetic networks: splits are minimal cuts. These model all the ways heredity can happen. Here, $N$ is 1-nested: each edge in at most 1 cycle.

Phylogenetic trees: every edge is a split of $[n]$.
$A|B=\{2,3\}|\{1,4,5,6,7,8\}$
2) Split networks: splits are parallel classes of edges. These model all the ways that genetic distances can be measured.
Here, $s$ is circular: it can be drawn on the plane with leaves on the exterior in circular order $c$.


## Ordering



Partial ordering on both sets: $x \leqslant y$ when the set of splits displayed by $x$ is contained in that displayed by $y$.
(Two networks, split or phylogenetic, are considered equivalent when their sets of splits are the same.)


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Mapping: [Gambette, Huber, Scholz, 2017]



Both maps are monotone.

## Galois connection: Reflection






$$
L(s) \leq N \Longleftrightarrow s \leq \Sigma(N)
$$

$L$ is surjective; $\Sigma$ is injective.

## Polytopes

## Definition

For a binary, level-1 phylogenetic network $N$, the vector $\mathbf{x}(N)$ is defined to have lexicographically ordered components $x_{i j}(N)$ for each unordered pair of distinct leaves $i, j \in[n]$ as follows:

$$
x_{i j}(N)= \begin{cases}2^{k-b_{i j}} & \text { if there exists } c \text { consistent with } N ; \text { with } i, j \text { adjacent in } c, \\ 0 & \text { otherwise. }\end{cases}
$$

where $k$ is the number of bridges in $N$ and $b_{i j}$ is the number of bridges crossed on any path from $i$ to $j$.
The convex hull of all the $\mathbf{x}(N)$ with $k$ nontrivial bridges is the level- 1 network polytope $\operatorname{BME}(n, k)$.


$(4,2,1,0,1,2,1,0,1,2,0,2,4,0,4)$
$(2,0,1,0,1,2,0,0,0,1,0,1,2,0,2)$

## Result: faces of $\operatorname{BME}(n, k)$

## Theorem

Every $n$ leaved 1-nested network $N$ corresponds to a face of each $\operatorname{BME}(n, k)$ polytope.

That face has vertices all the binary level-1 $k$-bridge networks $N^{\prime}$ such that $N \leq N^{\prime}$.




## Counting vertices



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$1 \cdot 3$

## Counting vertices



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$1 \cdot \frac{(n-1)!}{2}$

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$1 \cdot \frac{(n-1)!}{2} T(n, k)$

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$$
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$$

## Counting vertices





$1 \cdot \frac{(n-1)!}{2} T(n, k) \frac{1}{2^{k}}$

Counting vertices: note the cases $k=0, k=n-3$.


## Weighting

A weighting of a split network $s$ is a function of splits, $w: s \rightarrow \mathbb{R}_{\geq 0}$.
For a weighted split network $s$ we define the distance vector $\mathbf{d}_{s}$ where

$$
\mathbf{d}_{s}(i, j)=\sum_{i \in A, j \in B} w(A \mid B)
$$

where the sum is over all splits of $s$ with $i$ in one part and $j$ in the other.

A weighting for a phylogenetic network $N$ is a function from edges to positive reals.
We define $\mathbf{d}_{N}$ where

$$
\mathbf{d}_{N}(i, j)=\min \left\{\sum_{e \in P} w(e) \mid P \text { path } i \rightarrow j\right\}
$$

is the minimum sum of the weights of edges over all paths $i$ to $j$.

## Weighting

Theorem
For any weighted planar phylogenetic network $N$, the distance vector $\mathbf{d}_{N}$ obeys the Kalmanson condition: there exists a circular ordering $c$ of $[n]$ such that for all $1 \leq i<j<k<I \leq n$ in that ordering,
$\max \left\{\mathbf{d}_{N}(i, j)+\mathbf{d}_{N}(k, l), \mathbf{d}_{N}(j, k)+\mathbf{d}_{N}(i, l)\right\} \leq \mathbf{d}_{N}(i, k)+\mathbf{d}_{N}(j, l)$.

Theorem
Theorem: For any Kalmanson vector $\mathbf{d}$, there is a unique weighted split network $s=N N(\mathbf{d})$ such that $\mathbf{d}_{s}=\mathbf{d}$.

## Weighting

For any Kalmanson d for [n], we can restrict the poset of phylogenetic networks (with weighting) to those that have $\mathbf{d}_{N}=\mathbf{d}$.

For a weighted split network $s$ define $L_{w}(s)$ to be $L(s)$ with weight function

$$
w_{s}(e)=\sum_{e \in C(A \mid B)} w(A \mid B)
$$

where the sum is of weights in $s$ of splits whose cuts the edge is a member of. For a weighted phylogenetic network $N$ define

$$
S_{w}(N)=N N\left(d_{N}\right)
$$

## Galois connection: Coreflection






$$
\begin{gathered}
L_{w}(s) \leq N \Longleftrightarrow s \leq S_{w}(N) \\
L_{w} \text { is } 1-1 ; S_{w} \text { is onto. }
\end{gathered}
$$

## Result

By comparing the two connections and the face theorem, we see that:

Theorem
Given any weighted phylogenetic network $N$, the product $\mathbf{x}(\hat{N}) \cdot \mathbf{d}_{N}$ is minimized simultaneously for the binary networks $\hat{N}$ with $k$ bridges such that $\overline{S_{w}(N)} \leq \Sigma(\hat{N})$.
In $\overline{S_{w}(N)}$ the overline indicates forgetting the weighting.

## Polytope Pictures: Duality


$\operatorname{BME}(4,0)$ on left, with $\operatorname{BME}(4,1)$ included.
$\operatorname{CSN}(4)$ on the right, with $\mathrm{BHV}(4)$ included.

## Facets



## Facets



## Nesting



## Nesting



## Nesting



## Nesting



## Nesting



Thanks so much!

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Questions...

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What are the antipodes of the incidence Hopf algebras of these posets?

